

المقدمة التاسعة

File Systems

File system consist of:

1. Files.
 2. Directory structure, which organize all files in the system.
- **File:** A logical storage unit mapped by the O.S onto physical devices. It is a collection of related information recorded on nonvolatile secondary storage.

File Attributes

1. Name
 2. Type
 3. Location
 4. Size
 5. Protection: read only, executing, ...
 6. Time, date, and user identification.
- The information about all files is kept in the directory structure.

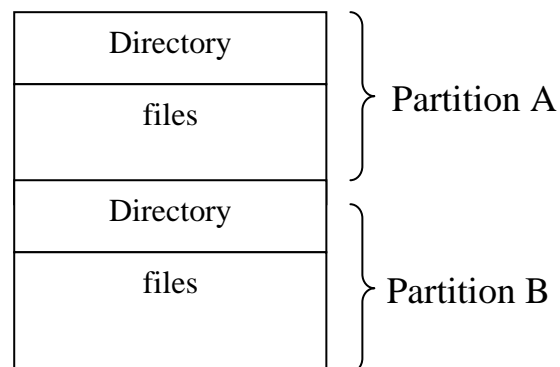
Access Methods

1. **Sequential access:** information in the file is processed in order, one record after the other. Ex: editors and compilers.
2. **Direct access:** reading or writing records rapidly in no particular order. Ex: data base.

Directory Structure

To manage all files in the system, we need to organize them. This organization is done in two parts:

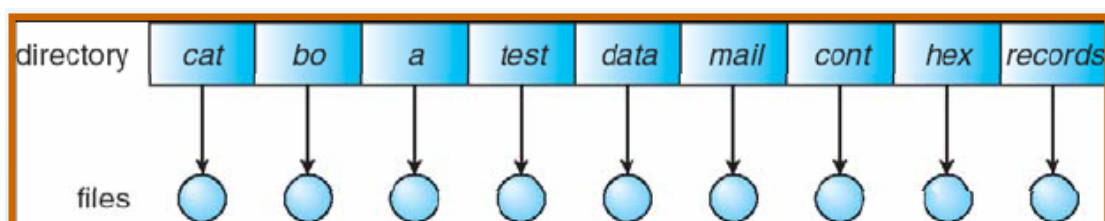
1. **Partitions:** file system is broken into partitions, to provide several separate areas within one disk.
2. **Directory:** each partition contain information about files within it in directory, which records information such as name, location, size, and type for all files on that partition.



Directory Schemes

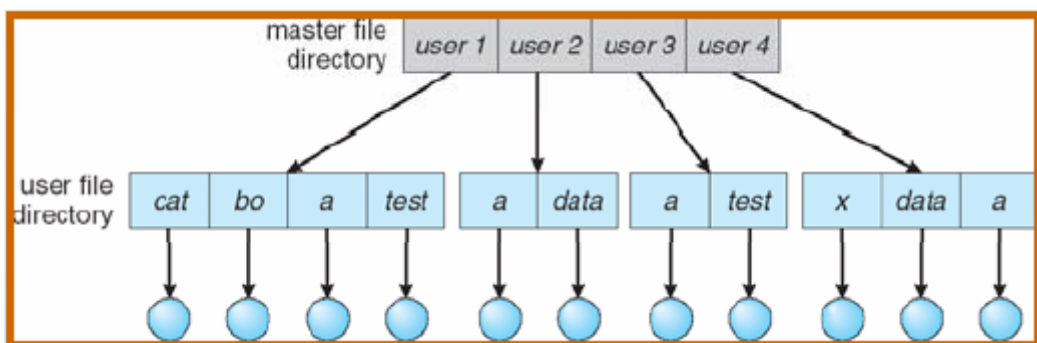
1. Single-level directory

- The simplest structure.
- All files are contained in the same directory.
- Disadvantage: confusion of file names of different users (***name –collision problem***).



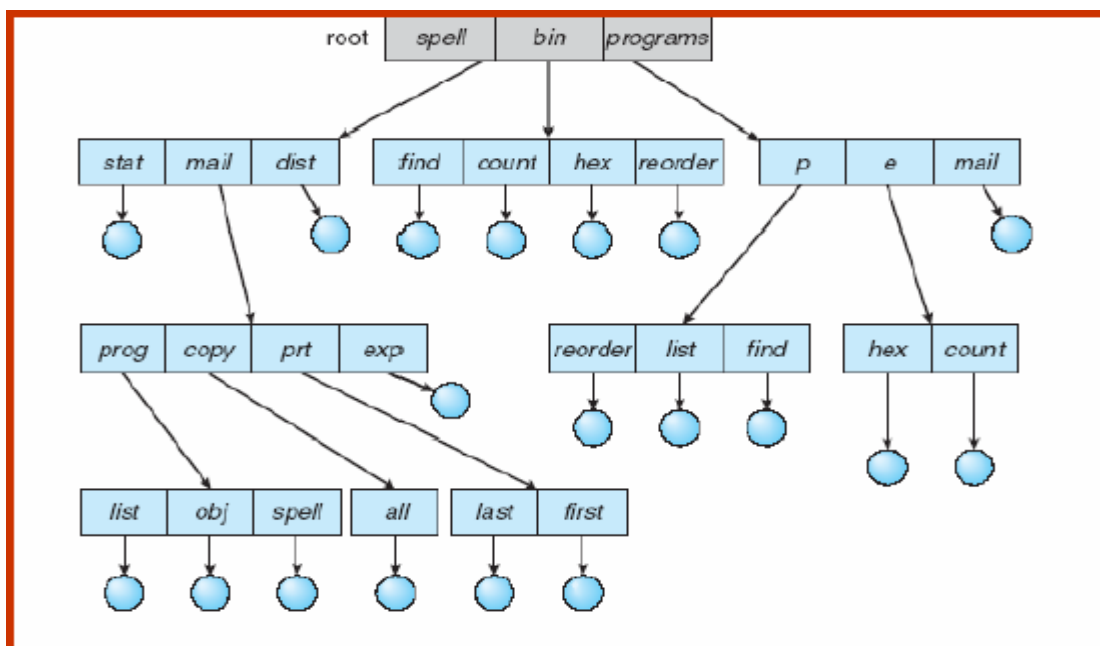
2. Two-level Directory

- Create a separate directory for each user.
- When a user job starts a user logs in the System's Master File directory (MFD) is searched by the user name for User File Directory (UFD).
- This scheme solve the name-collision problem.
- Its disadvantage when users want to cooperate on same task and to access one another's files.



3. Tree-Structure Directories

- This scheme allows user to create subdirectories and files.
- Subdirectory contains a set of files or subdirectories.
- Every file has a unique path name.

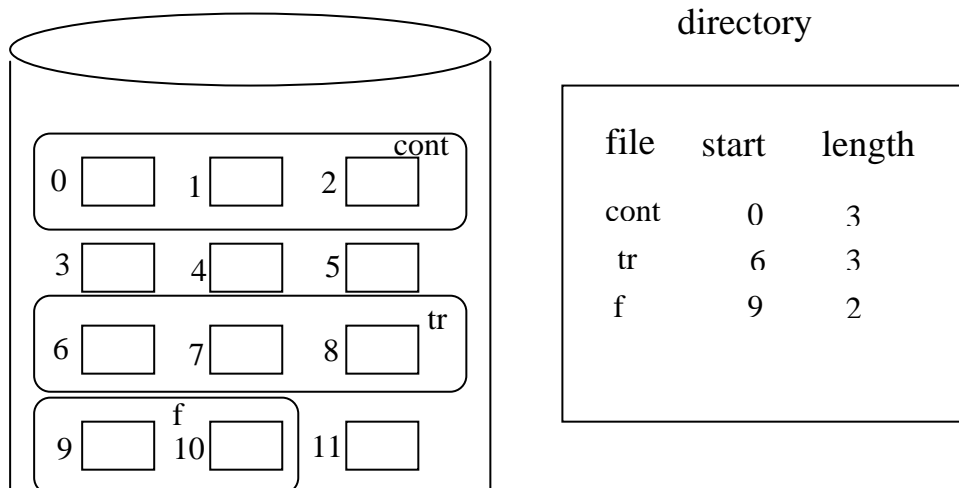


Allocation Methods

How to allocate space to files? so that disk space is utilized effectively and files accessed quickly.

1. Contiguous Allocation:

- Each file should occupy a set of contiguous blocks on the disk.
- A directory entry for each file indicates address of a starting block and length of area of this file.

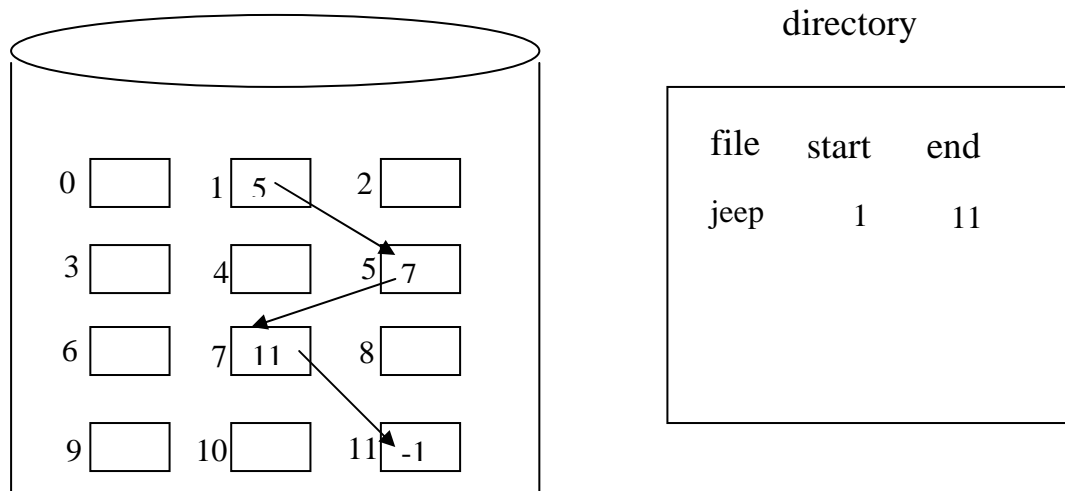


- Suffer from external and internal fragmentation.
- Difficult to find space for a new file and extend the file.

2. Linked Allocation:

Each file is a linked list of disk blocks. These blocks may be scattered anywhere on the disk.

- The directory have a pointer to the first and last blocks of the file.



Advantages:

- a. no external fragmentation.
- b. No need to declare the file size.
- c. No necessary for compaction.

Disadvantages:

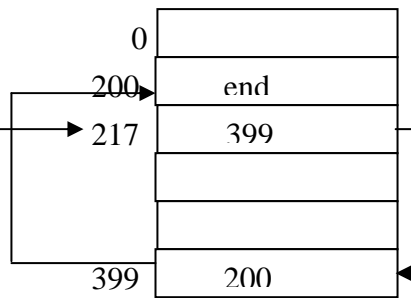
- a. used for only sequential accessed files.
 - b. Space required for pointers.
- Solution for pointer space problem is to collect blocks into (**clusters**) and allocate the cluster rather than blocks.
 - An important variation on the linked allocation method is the use of **FAT (File Allocation Table)**. Which is simple and efficient method of disk-space allocation.
 - it found in the beginning of each partition.
 - FAT is used as a linked list.
 - Directory entry contains the first block no.

Directory entry

test	217
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name

Start block

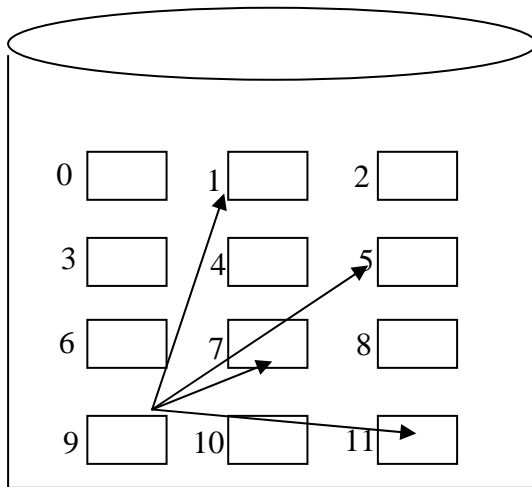


FAT

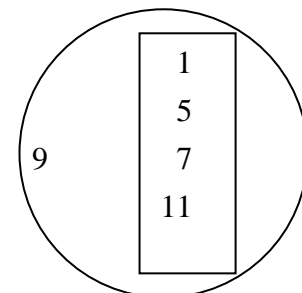
3. Indexed Allocation:

- Bring all the pointers together into one location : **index block.**
- The directory contains the address of the index block.

directory



file	Index block
jeep	9



- This organization support direct access without suffering from external fragmentation.