Effectiveness of the Limited Retransmission on the WLANs Using Error-Prone Channel

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Abstract

IEEE 802.11 based wireless local area networks (WLANs) is one of the most popular protocols for interconnecting portable communication devices. All the stations communicating over WLANs share a common wireless channel. Many conscious studies have been devoted to utilize this precious medium efficiently. However, most of these studies have been done either under assumption of idealistic channel condition or with unlimited retransmitting number. This paper is devoted to investigate influence of limited retransmissions and error level in the utilizing channel on the network throughput, probability of packet dropping and time to drop a packet. The results show that the networks using four-way handshaking mechanism has a good immunity against the error over the available range of retry limits, also the throughput is unchangeable with size of the network over the range of retransmission limits. On the other side, for the networks using basic access mechanism, the throughput is suppressed with increasing amount of errors in the transmitting channel over all the range of the retry limit as well as it is quite sensitive to the size of the network. However, the throughput does not change with retry limits when it exceeds the maximum number of the backoff stage in both DCF's mechanisms. In the both mechanisms the probability of dropping a packet is a decreasing function with number of retransmissions and the time to drop a packet in the queue of a station is a strong function to the number of retry limit, size of the network, the utilizing medium access mechanism and amount of errors in the channel.

Keywords*: IEEE802.11 DCF, WLAN, MAC protocol, Throughput, Error-prone channel.*

1. Introduction

One key point to satisfy most of demands within the next generation communications (subscribe accessible at any time and at anywhere) is to increase both data rate and speed [1, 2]. The Wireless local area networks (WLANs) are able to comply with such demands and have become one of the fastest growing segments in the communications industry. The worldwide shipments of the WLAN equipment products arrive $5.2 billion in 2005, and it is expected that WLAN equipment will continue growing in 2006 to reach around $5.9 billion level as new IEEE 802.11n and VoWi-Fi equipment is introduced and the infrastructure for traditional Wi-Fi expands [3, 4]. In 1997 IEEE's committee standardized 802.11 protocol for WLANs [5]. Since that time several versions of this protocol have been made. The physical media in the WLANs is shared between all stations and has limited connection range compared with its wired counterpart. The standard defines three PHY technologies and a unified MAC protocol to support 1 and 2 Mbps transmission over wireless media. The MAC protocol has two functions, namely distributed coordination function (DCF) and the optional point coordination function (PCF). DCF has superior attractiveness over PCF in many aspects [6], therefore this study is conducted to investigate WLANs utilizing DCF. DCF defines two mechanisms to access transmission medium: the basic access scheme, which is the default scheme and the request to send/clear to send (RTS/CTS) scheme, also known as four-way handshaking scheme [4, 7, 8]. Recently, considerable studies have been concentrated on modeling the IEEE 802.11 DCF medium access method. Bianchi in [9] modeled the idealistic assumption of collision only errors, that packet retransmits are unlimited and a packet is being transmitted continuously until its successful reception. Wu in [10] extended Bianchi's analysis to include the finite packet retry limits as defined in the standard. Both studies used Markov chain model to analyze DCF operation and calculated the saturated throughput of 802.11 protocol. Periklais et al. [11] extended the work in [9] and [10] by taking into account both: transmission errors and packet retry limits for basic access of the IEEE802.11a protocol. X. Wang et al. [12] evaluate the impact of transmission error rate on the contention and the system throughput in WLAN's protocol. However, [11] and [12] considered the probability of bit errors appearing on the transmission channel is the same in the two access mechanisms. Z. Tang et al. [13] presented an analytical model to evaluate the performance of the DCF in the case of bit errors appearing on the transmission channel and taking type of the used mechanism into account. However, Tang's study was under assuming of unlimited retransmitting. In this paper we extend Tang's work by taking into account influence of retransmissions and investigate its impacts on the performance of the WLANs. The results show that the throughput is insensitive to the number of the retransmissions when it exceeds the maximum number of

backoff stages in both mechanisms. This insensitiveness' "trend" does not change with amount of BER in the utilizing channel. However, adopting four-way handshaking mechanism show that the throughput is more immunes than its counterpart mechanism when WLAN's channel suffers from much error. This paper is organized as follows: Section 2 devoted to explain the used model in this study. Explanations for the achieved results are given in Section 3 and then our conclusions are drew to Section 4.

2. The Model

The analysis employs the Markov chain model in [8] and [10] and makes use of the same assumptions as in [13]: all stations always have a packet available for transmitting (saturation case) into an error prone-channel.

2.1. A brief description of the backoff process in IEEE 802.11 DCF

When stations sensing the medium is idle for a period, more than DCF, the backoff timer value for each station is uniformly chosen within the interval $[0, W_i - 1]$, where W_i is the current contention window (CW) size and i is the backoff stage $i \in [0, m]$ and m represents the station's retry limit. The backoff counter for every station depends on the collision and on the successful packet transmissions experienced by the station in the past. At the beginning: $W = W_o$ and after each retransmitting due to a packet collision or error, W_i is doubled up to a maximum value, $W_{m'} = 2^{m'} W$, where m' is the maximum number of the backoff stages. When W_i reaches $W_{m'}$, it will stay at this value until it reset to W_o again either after the successful data packet transmitting or when the counter reaches to its limit.

2.2. The analytical modeling[1]

As in [10], the probability of a station to transmit a packet in a randomly chosen slot time is:

$$\tau = \frac{1 - p^{m+1}}{(1 - p)} . \Psi_{0,0} \qquad (1)$$

where:

$$\Psi_{0,0} = \begin{cases} \dfrac{2(1-2p).(1-p)}{W(1-(2p)^{m+1}.(1-p)+(1-2p).(1-p^{m+1}))} & m \le m' \\[4mm] \dfrac{2(1-2p).(1-p)}{W(1-(2p)^{m+1}).(1-p)+(1-2p).(1-p^{m+1})+W.2^{m'}p^{m'+1}.(1-2p).(1-p^{m-m'})} & m > m' \end{cases}$$

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[1] Definition and values of all the rest parameters that do not mentioned here are as in Ref. [13].

τ does not depend on the type of the mechanism adapted by a station: basic access or four-way handshaking, p is the unsuccessful probability when a transmitted packet encounters a collision with at least one of the $n-1$ remaining stations in a time slot. So:

$$p = 1 - (1-\tau)^{n-1}(1-p_c) \qquad (2)$$

Influence of errors in the transmitting channel included through the parameter p_c as [13]: In the case of basic access mechanism:

$$p_c = 1 - (1 - BER)^L \qquad (3\text{-}a)$$

where $L = PHY_h + MAC_h + P + ACK$. In the case of four-way handshaking:

$$P_c = 1 - (1-BER)^{RTS} * (1-BER)^{CTS} * (1-BER)^P * (1-BER)^{ACK} \qquad (3\text{-}b)$$

When a station transmits and the remaining $n-1$ stations defer their transmissions, the packet would be arrive successfully with probability p_s. Considering $(1-p_{tr})$ as the probability of a random slot is empty, and probability of successful transmission is $p_{tr} p_s$ and probability of the collision as $p_{tr}(1-p_s)$, the average length of a slot time is:

$$E[slot] = (1 - p_{tr}).\sigma + p_{tr}.p_s.T_s + p_{tr}.(1-p_s).T_c \qquad (4)$$

Consequently, the system throughput, S, can be expressed by dividing the successfully transmitted payload data over a slot time. The probability of dropping a packet when the retry limit is reached is known as the packet drop probability and given as:

$$p_{drop} = p^{m+1} \qquad (5)$$

A packet is dropped when it reaches the last backoff stage and experiences another collision or an error.

3- Performance analysis

This analysis is based on the model in [13] and we have included influence of limited number of retransmissions. To validate our evaluation and highlight influence of limited retries we compare our results, *after* taking influence of limited retries, with results have been gotten by using the model in [12] and with that using model given in Ref. [13]. The results are illustrated in Fig. 1. The same parameters values in [13] have been used in this study in order to facilitate the comparison purpose. In Fig. 1 our results denoted as (a), results of Ref. [13] denoted as (b) and the results that obtained by using the model in [12] denoted as (c). The system throughput has been estimated for three different network sizes: 5 (small), 20 (middle), and 50 (large). Fig. 1 shows that: for networks utilizing the basic access, results (a) are much closer to results (b) in the small and middle network sizes and much close to results (c) for the large networks over *all* the range of BERs. This can be justify since on the small and middle networks the rate of

collisions is relatively small compared with that in large networks, which consequences influence of retransmission at small and middle networks is also small. For networks utilizing the four-way handshaking, all results: (a), (b) and (c) show that the throughputs are insensitive to size of the networks over all the examined range of the BERs. While almost all the results are mostly closed when the BERs less than 10^{-5} there is a difference between the results (a) and (b) in one side and results (c) on the other side when $10^{-5} < BER < 5 \times 10^{-4}$. This difference is due to the effect of the parameter p_c on p, which could be much sensible at

higher level of BERs. Considering errors in the channel, p_c on the Ref. [12] did not distinguish the used access mechanism as that done in [13] and considered here as shown in Eq. 3.

Figure 2- a, and b illustrate variation of the network throughput with number of retransmissions at low level of BERs (= 10^{-5}) as in Fig. 2- a and at relatively high level of BERs (= 10^{-4}) as stated in Fig. 2-b. It is obvious that trend of the network throughput is almost unchangeable with level of errors in the channel when a network uses the four-way handshaking scheme. In the contrary side, when the networks use basic access, the throughput level is obviously decreased with increasing errors in the channel over all the retransmission range. Fig. 3-a and b show that probability to drop a packet is exponentially decreasing with number of retransmissions and has undistinguishable differences between the two access mechanisms. However, this probability is increasing function of errors in the channel. Probability of dropping a packet is ignorable when the network exceeds the maximum number of backoff stage[2] with low level of BER (see Fig. 3 a). This leads to enforce packets stay at queue of the transmitter for a long time and cause much delay, which could be unacceptable for some sensitive time-delay applications. Furthermore, Fig. 4 a and b show that the average time to drop a packet when adapting four-way handshaking is lower than that when using basic access at low level of BER that is obvious at high level of BER (as in Fig 4-b). These are beneficial of using RTS/CTS packets by reserving the channel in advance to start transmitting a long data packet and hence supporting reduction the rate of collisions.

4- Conclusions:

This paper presents an analytical model included influence of limited number of retransmissions on the main characteristics of WLANs that use error-prone channels. The model has been applied on the two available mechanisms in the DCF functions. We validate our

evaluations by comparing our results with the results in two considered references. This study shows that: the networks using four-way handshaking mechanism have a good immunity against increasing the error in the transmitting channel over the range of retransmissions, also WLAN's throughput is unchangeable with size of the network over the range. On the other side, for the networks using basic access mechanism, the throughput is suppressed with increasing amount of errors in the transmission channel over *all* the available range of the retry limits as well as it is sensitive to the size of the network. Further, the throughput does not change with retry limits when it exceeds the maximum number of backoff stages in both DCF's mechanisms. In the both mechanisms probability of dropping a packet is decreasing function with number of retransmissions and the time to drop a packet in the queue of a station is strong function to the number of retry limits, size of the network, the utilized mechanism and the amount of errors in the transmitting channel.

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[2] Extensive calculations for different values have been done, and all the obtained results were valid and assist the result stated in this paper: Probability of dropping a packet is ignorable when the network exceeds the maximum number of backoff stage.

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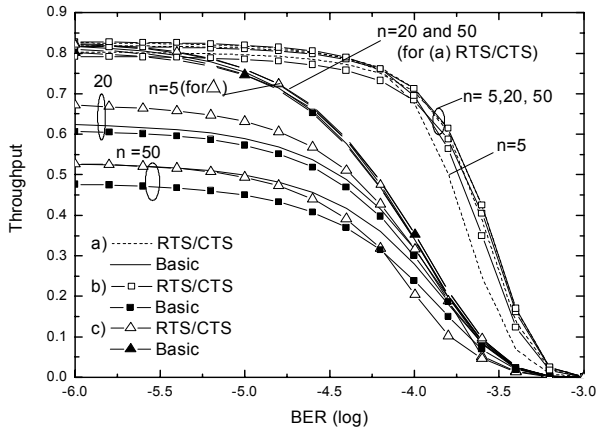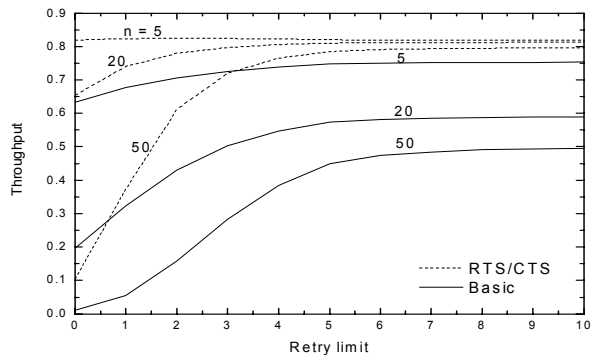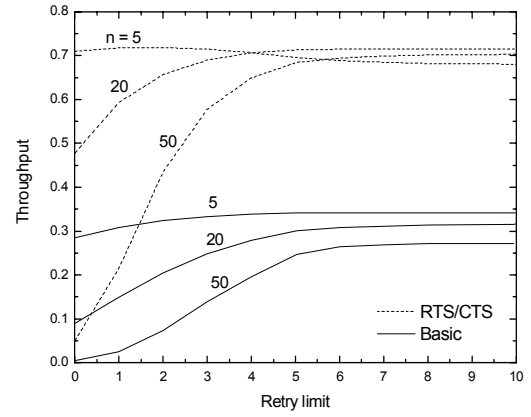
Figure 1 Throughput comparison for different network size (n=5, 20 and 50): (a) our results, (b) results from Ref. [13], and (c) results obtained using the model in Ref. [12]
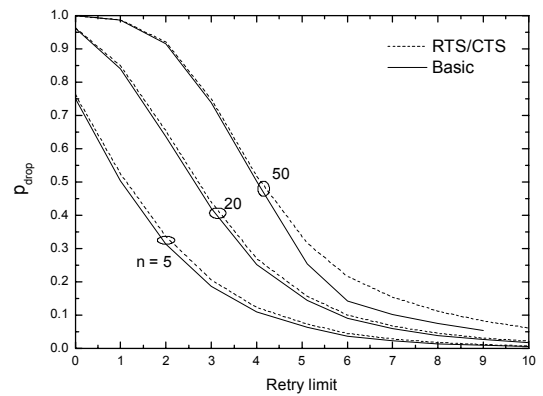
(a) BER = 10^{-5}

(b) BER = 10^{-4} .

Figure 2 Throughput efficiency of the two mechanisms as a function of retry limit for different network sizes
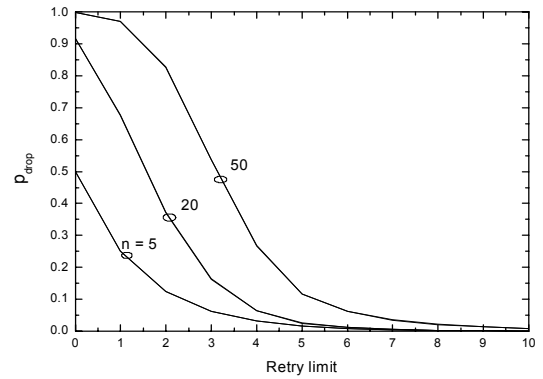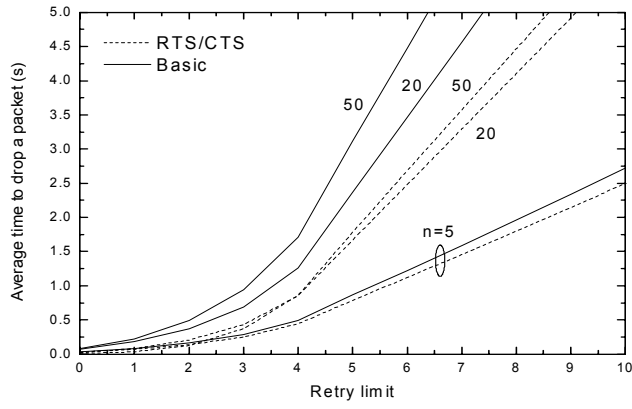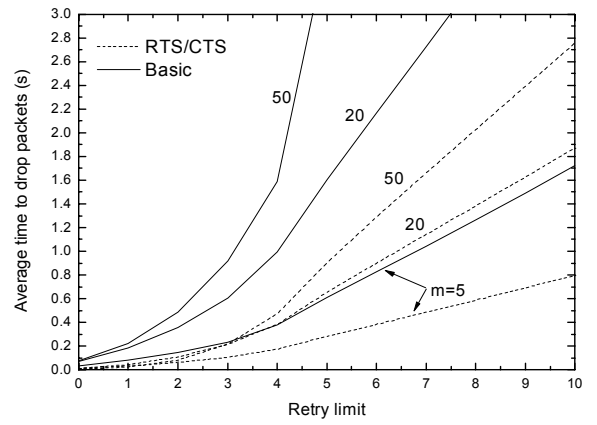
(a) BER = 10^{-5}

(b) BER = 10^{-4}

Figure 3 Impact of retry limit on the drop probability for different network sizes

(a) BER = 10^{-5}

(b) BER = 10^{-4}

Figure 4 Influence of retry limit on the time to drop a packet for different network sizes。