

Edge addition and edge deletion of graphs^{*}

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Abstract: Let $P(t, d)$ (resp. $C(t, d)$) be the minimum diameter of a connected graph obtained from a single path (resp. from a single circle) of length d by adding t extra edges. It is proved that $P(t, d) = \left\lceil \frac{d-2}{t+1} \right\rceil + 1$ if t and d satisfy the following conditions: $t \geq 4$ and $t+4 \leq d \leq t+7$, $t=4$ and $d=10k+1$ ($k \geq 1$). For some t and d , the exact value and the best lower bound of $C(t, d)$ are determined, and so the conjecture of Schoone, et al [J. Graph Theory, 1987, 11:409-427] is settled partially.

Key words: diameter; altered graph; edge addition; edge deletion; Schoone et al's conjecture

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图的边添加和减少

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摘要: 用 $P(t, d)$ (或者 $C(t, d)$) 表示从一条长为 d 的简单路(或者简单圈)通过添加 t 条边后得到图的最小直径. 证明了: 如果 t 和 d 满足条件 $t \geq 4$ 且 $t+4 \leq d \leq t+7$, 或者 $t=4$ 且 $d=10k+1$ ($k \geq 1$), 那么 $P(t, d) = \left\lceil \frac{d-2}{t+1} \right\rceil + 1$. 对某些 t 和 d , 确定了 $C(t, d)$ 的值和最好下界, 部分地解决了 Schoone 等的猜想[J. Graph Theory, 1987, 11:409-427].

关键词: 直径; 变更图; 边添加; 边减少; Schoone 等的猜想

0 Introduction

We follow ref. [1] for graph-theoretical terminology and notation not defined here. Let G be a simple undirected graph with vertex-set $V = V(G)$ and edge-set $E = E(G)$.

Let $P(t, d)$ denote the minimum diameter of a graph obtained by adding t extra edges to a path

$P = (x_1, x_2, \dots, x_{d+1})$ of length d . For some small t 's and special d 's, the values of $P(t, d)$ have been determined. It is easy to verify that $P(1, d) = \left\lceil \frac{d+1}{2} \right\rceil$ for $d \geq 2$. Ref. [2] determines $P(2, d) = \left\lceil \frac{d+1}{2} \right\rceil$ for $d \geq 3$ and $P(3, d) = \left\lceil \frac{d+2}{4} \right\rceil$ for $d \geq 5$.

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Ref. [3] determines $P(t, (2k-1)(t+1)+1) = 2k$ for any positive integer k ,

$$\left\lceil \frac{d}{t+1} \right\rceil \leq P(t, d) \leq \left\lceil \frac{d}{t+1} \right\rceil + 1$$

for $t = 4, 5$ and $d \geq 4$ and in general

$$\left\lceil \frac{d}{t+1} \right\rceil \leq P(t, d) \leq \left\lfloor \frac{d-2}{t+1} \right\rfloor + 3.$$

Also ref. [4] determines

$$P(t, d) \leq \begin{cases} \left\lfloor \frac{d-2}{t+1} \right\rfloor + 2 & \text{if } d \in I'(t, k), \\ \left\lceil \frac{d-2}{t+1} \right\rceil + 1 & \text{otherwise,} \end{cases}$$

where $I'(t, k) = \{2k(t+1)+1, 2k(t+1)+2, 2k(t+1)-t+1\} \cup \{2k(t+1)-t+h; h=6, 7, \dots, t\}$ and the integers $t \geq 6, d \geq 2$ and $k \geq 1$.

Let $C(t, d)$ denote the minimum diameter of a graph obtained by adding t extra edges to a circle of length d . It is easy to verify that $C(1, d) = \left\lfloor \frac{d}{2} \right\rfloor$ for $d \geq 2$. Ref. [2] determines $C(2, d) = \left\lceil \frac{d+2}{4} \right\rceil$ for $d \geq 4$. For general $t \geq 3$ and $d \geq 5$, ref. [5] obtains $\left\lceil \frac{d}{t+i} \right\rceil - 1 < C(t, d) \leq \left\lceil \frac{d}{t+i} \right\rceil + 3$ for i equal 1 if t is odd and 2 if it is even. WU and XU^① prove that for $d \geq 5, \left\lceil \frac{d}{4} \right\rceil - 1 < C(3, d) \leq \left\lceil \frac{d}{4} \right\rceil$ and in general $C(t, d) \leq \left\lceil \frac{d-9+i}{t+i} \right\rceil + 3$ for i equal 1 if t is odd and 2 if t is even for any integer $t \geq 4$.

Let $f(t, d)$ denote the maximum diameter of a connected graph obtained after deleting t edges from a connected graph of diameter d . For $t = 1$, ref. [6] determines $f(1, d) = 2d$. For small t and d , Schoone, et al^[2] prove that $f(2, d) = 3d - 1$, $f(3, d) = 4d - 2$ for $d > 1$, $f(t, 2)$ is equal to $t + 3$ for $t = 1, 2, 3, 4, 6$ and $t + 2$ otherwise, $f(t, d) \leq (t+1)d$ for any positive integers t and d , and $f(t, d) \geq (t+1)d - t$ for d is even. Also WU and XU^① prove $f(t, d) \geq (t+1)d - 2t + 4$ for $t \geq 4$ and d is odd not less than 3.

In this paper we first prove that $P(t, d) =$

$\left\lceil \frac{d-2}{t+1} \right\rceil + 1$, for $t \geq 4, t+4 \leq d \leq t+7$ and for $t = 4, d = 10k+1$ ($k \geq 1$). Then we prove that for $t \geq 3, C(t, d) = 3$ for $t+6 \leq d \leq t+8$, and the lower bound $C(t, d)$ is 3 for $d = t+9$ and $2k$ for $d = (2k-1)(t+2)+2$ ($k \geq 1$). These bounds are tight for some t and d .

After that we prove the conjecture of Schoone et al^[2] $f(t, d) \leq (t+1)d - t + 1$, for $t \geq 4, t+4 \leq d \leq t+7$, and $d = 2k(t+1) - t, k \geq 1$, and for $t = 4$ and $d = 10k + 1$. This upper bound is tight for some d .

1 Several lemmas

Lemma 1.1^[4] For any integer $k \geq 1$, let $I'(t, k) = \{2k(t+1)+1, 2k(t+1)+2, 2k(t+1)-t+1\} \cup \{2k(t+1)-t+h; h=6, 7, \dots, t\}$. Then

$$P(t, d) \leq \begin{cases} \left\lfloor \frac{d-2}{t+1} \right\rfloor + 2 & \text{if } d \in I'(t, k), \\ \left\lceil \frac{d-2}{t+1} \right\rceil + 1 & \text{otherwise,} \end{cases}$$

for the integers $t \geq 6$, and $d \geq 2$.

Lemma 1.2^[3] For given positive integers t and d ($d \geq 2$), $\left\lceil \frac{d}{t+1} \right\rceil \leq P(t, d) \leq \left\lfloor \frac{d-2}{t+1} \right\rfloor + 3$. In particular, $P(t, (2k-1)(t+1)+1) = 2k$ for any positive integer k , $P(t, d) \leq \left\lfloor \frac{d}{t} \right\rfloor + 1$ if k is large enough, and $\left\lceil \frac{d}{t+1} \right\rceil \leq P(t, d) \leq \left\lceil \frac{d}{t+1} \right\rceil + 1$ for $t = 4, 5$ and $d \geq 4$.

Lemma 1.3^[7] Let G be a connected undirected graph, $S \subset E(G)$ with $|S| = t$. If $h = d(G-S)$ is well defined, then $d(G) \geq P(t, h)$.

Lemma 1.4^① For $d \geq 5$ and $t \geq 4$, $C(t, d) \leq$

$$\begin{cases} \left\lfloor \frac{d-7}{t+2} \right\rfloor + 3 & \text{if } t \text{ is even,} \\ \left\lceil \frac{d+t-6}{2t+2} \right\rceil + \left\lceil \frac{d+t-1}{2t+2} \right\rceil \leq \frac{d-8}{t+2} + 3 & \text{if } t \text{ is odd.} \end{cases}$$

Also $\left\lceil \frac{d}{4} \right\rceil - 1 \leq C(3, d) \leq \left\lceil \frac{d}{4} \right\rceil$.

① WU Ye-zhou, XU Jun-ming. On diameters of altered graph, to appear in J. Math. Research Exposition (2006).

2 Proof of main results

2.1 Edge addition

Theorem 2.1 For $t \geq 4$,

$$P(t, d) = \left\lceil \frac{d-2}{t+1} \right\rceil + 1 = 3,$$

where $t+4 \leq d \leq t+7$.

Proof Let $P = (x_0, x_1, \dots, x_d)$ be an (x_0, x_d) -path and G an altered graph obtained from P by adding t extra edges and having diameter $d(G) = P(t, d)$, where $d \geq t+4$.

If $d(G) = 1$, then the number of the extra edges is equal to

$$t \geq (d-1) + (d-2) + \dots + 1 = d(d-1)/2 \geq d-3 = t+1,$$

a contradiction.

If $d(G) = 2$, then let x_i be the smallest numbered vertex that G has no edge (x_i, x_j) with $j > i+1$. For each $j = 0, 1, \dots, i-1$, there exists a j' ($j' \geq j+2$) such that $(x_j, x_{j'}) \in E(G)$ is an extra edge, and so such edges are at least i .

Since $d(G) = 2$, we must be able to reach every other vertex in two steps from x_i . Hence we need edges $(x_{j'}, x_j)$ with $j \geq j'+2$ for all j with $i+3 \leq j \leq d$, since these vertices could not be reached in two steps from x_i in P . Such edges are at least $d+1-(i+3)-1 = d-i-3$, where the appearance of (-1) is due to the possibility that there may be an extra edge (x_{i-1}, x_j) in the first (i) extra edges when $j' = i-1$. Note that if $j' < i-1$, then $(x_i, x_{j'})$ and $(x_{j'}, x_j)$ are extra edges. Thus

$$t \geq i + (d-i-3) = d-3 = t+1,$$

a contradiction.

Thus $P(t, t+4) = d(G) \geq 3$. Since $P(t, d) \leq P(t, d')$ if $d \leq d'$, $P(t, d) \geq 3$ for $d \geq t+4$. On the other hand, since $t+4, t+5, t+6, t+7 \notin I'(t, k)$ for $k = 1$, from Lemma 1.1,

$$P(t, d) \leq \left\lceil \frac{d-2}{t+1} \right\rceil + 1 = 3 \text{ for } t \geq 6.$$

For $t = 4, 5$, from Lemma 1.2, $P(t, d) \leq \left\lceil \frac{d}{t+1} \right\rceil + 1 = 3$ for $d \geq 4$. Thus, $P(t, d) = 3$ for $t+4 \leq d \leq t+7$ and $t \geq 4$. \square

Theorem 2.2 $P(4, d) = \left\lceil \frac{d-2}{5} \right\rceil + 1$ where d

$= 10k+1$ and $k \geq 1$.

Proof First we prove that $P(4, d) \leq \left\lceil \frac{d-2}{5} \right\rceil + 1$, where $d = 10k+1$ and $k \geq 1$. We add four edges

$$e_1 = x_1 x_{4k}, \quad e_2 = x_{4k} x_{8k+2},$$

$$e_3 = x_{2k} x_{6k+1}, \quad e_4 = x_{6k+1} x_{d+1}$$

to the path $P = x_1 x_2 \dots x_{d+1}$ of length d (see Fig. 1 for $k = 2$ and $d = 21$). Now the end-vertices of these edges divide P into five segments

$$\begin{aligned} L_1 &= P(x_1, x_{2k}), & L_2 &= P(x_{2k}, x_{4k}), \\ L_3 &= P(x_{4k}, x_{6k+1}), & L_4 &= P(x_{6k+1}, x_{8k+2}), \\ L_5 &= P(x_{8k+2}, x_{d+1}). \end{aligned}$$

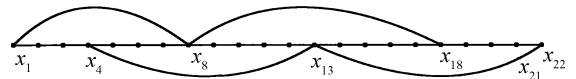


Fig. 1 Construction of Theorem 2.2 for $k = 2$ and $d = 21$

Define ten cycles as follows

$$\begin{aligned} C^1 &= L_1 \cup L_2 + e_1, \\ C^2 &= L_1 \cup L_3 + e_1 + e_3, \\ C^3 &= L_1 \cup L_4 + e_1 + e_2 + e_3, \\ C^4 &= L_1 \cup L_5 + e_1 + e_2 + e_3 + e_4, \\ C^5 &= L_2 \cup L_3 + e_3, \\ C^6 &= L_2 \cup L_4 + e_2 + e_3, \\ C^7 &= L_2 \cup L_5 + e_2 + e_3 + e_4, \\ C^8 &= L_3 \cup L_4 + e_2, \\ C^9 &= L_3 \cup L_5 + e_2 + e_5, \\ C^{10} &= L_4 \cup L_5 + e_4. \end{aligned}$$

Their lengths are

$$\begin{aligned} \epsilon(C^1) &= 4k; \\ \epsilon(C^i) &= 4k+2, \quad \text{for } i = 2, 5; \\ \epsilon(C^{10}) &\leq 4k+1; \\ \epsilon(C^i) &\leq 4k+3; \quad \text{for } i = 3, 4, 6, 7, 8, 9. \end{aligned}$$

Then any vertices x and y of G are contained in some cycles C^i defined above. Thus

$$\max\{d(C^i) : 1 \leq i \leq 10\} \leq \left\lceil \frac{4k+3}{2} \right\rceil = 2k+1.$$

This means that for $d = 10k+1$ and $k \geq 1$,

$$P(4, d) \leq d(G) \leq 2k+1 = \left\lceil \frac{d-2}{5} \right\rceil + 1.$$

From Lemma 1.2 we have

$$\left\lceil \frac{10k+1}{5} \right\rceil \leq P(4, 10k+1).$$

Since

$$\left\lceil \frac{10k+1}{5} \right\rceil = \left\lceil \frac{10k-1}{5} \right\rceil + 1$$

for $k \geq 1$, the theorem follows. \square

Theorem 2.3 For $t \geq 3, C(t, d) = 3$ where $t+6 \leq d \leq t+8; 3 \leq C(t, d) \leq 4$ for $d = t+9$; and $2k \leq C(t, d) \leq 2k+1$ for $d = (2k-1)(t+2) + 2$ and t is even and $k \geq 1$ or t is odd and $k = 1, 2, 3$.

Proof It is easy to verify that,

$$C(t, d+1) \geq P(t+1, d),$$

since one way of adding $t+1$ edges to a path P_{d+1} is to first add one edge joining two end vertices of P_{d+1} and then to add t edges in an optimal way to result in a cycle C_{d+1} . Then from Theorem 2.1 we have

$$C(t, d+1) \geq 3 \quad \text{for } t+5 \leq d \leq t+8, t \geq 3.$$

Also from Lemma 1.2 we have

$$C(t, d+1) \geq 2k, \text{ for } t \geq 3,$$

$$d = (2k-1)(t+2) + 1, k \geq 1.$$

Now from Lemma 1.4 we have for $t \geq 3$

$$C(t, d+1) \leq 3, \quad \text{for } t+5 \leq d \leq t+7,$$

$$C(t, d+1) \leq 4, \quad \text{for } d = t+8,$$

and $C(t, d+1) \leq 2k+1$ for $d = (2k-1)(t+2) + 1$, and t is even and $k \geq 1$ or t is odd and $k = 1, 2, 3$.

So the theorem follows. \square

2.2 Edge deletion

Theorem 2.4 $f(t, d) \leq (t+1)d - t + 1$ for $t \geq 4$ and $t+4 \leq d \leq t+7$, and $d = 2k(t+1) - t$ and $k \geq 1$, also for $t = 4$ and $d = 10k + 1$.

Proof Let $t \geq 4$. From Theorem 2.1 we have

$$P(t, d) = \left\lceil \frac{d-2}{t+1} \right\rceil + 1, \quad \text{for } t+4 \leq d \leq t+7.$$

From Theorem 2.2 we have

$$P(4, d) = \left\lceil \frac{d-2}{5} \right\rceil + 1, \text{ for } d = 10k + 1 \text{ and } k \geq 1.$$

From Lemma 1.2 we have

$$P(t, (2k-1)(t+1) + 1) = 2k.$$

This means

$$P(t, d) = \left\lceil \frac{d-2}{t+1} \right\rceil + 1,$$

$$\text{for } d = (2k-1)(t+1) + 1, k \geq 1.$$

Now let G be an undirected graph with diameter $d, S \subset E(G)$ with $|S| = t$ such that $d(G-S) = h = f(t, d)$. Thus from Lemma 1.3 and for $t \geq 4, d \in I_2(t, 1)$, and $d = 2k(t+1) - t + 5$, also for $t = 4$ and $d = 10k + 1$ and $k \geq 1$, we have

$$\left\lceil \frac{h-2}{t+1} \right\rceil + 1 = P(t, h) \leq d.$$

Then

$$\frac{h+t-1}{t+1} \leq d.$$

Thus

$$f(t, d) = h \leq (t+1)d - t + 1.$$

The theorem follows. \square

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